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演習第三

1. Let $w \in \{0, 1\}^*$ be any given string. We define

$$L_w = \{u w v \mid u, v \in \{0, 1\}^*\},$$

i.e., L_w is the set of all strings over $\{0, 1\}$ containing w as a *substring*.

- (1) Construct a nondeterministic finite automaton $\mathcal{A} = [\{0, 1\}, Q, \delta, q_0, F]$ accepting L_{00100} .
- (2) Transform \mathcal{A} into a deterministic finite automaton that also accepts L_{00100} .

2. Consider the following nondeterministic finite automaton $\mathcal{A} = [\{0, 1\}, Q, \delta, q_0, F]$, where $Q = \{q_0, q_1, q_2, q_3, q_4, q_5, q_6, q_7\}$ and $F = \{q_7\}$, and

$$\begin{aligned} \delta(q_0, 0) &= \{q_0, q_1\}, & \delta(q_0, 1) &= \{q_0, q_2\}, & \delta(q_1, 0) &= \{q_3\}, & \delta(q_2, 1) &= \{q_4\} \\ \delta(q_3, 0) &= \{q_5\}, & \delta(q_4, 0) &= \{q_6\}, & \delta(q_5, 1) &= \{q_7\}, & \delta(q_6, 0) &= \{q_7\} \end{aligned}$$

- (1) Construct a deterministic finite automaton \mathcal{A}' accepting the same language as \mathcal{A} .
 - (2) Draw the state diagrams of \mathcal{A} and \mathcal{A}' , respectively.
 - (3) Construct a regular grammar \mathcal{G} such that $L(\mathcal{G}) = L(\mathcal{A})$.
3. Let $\Sigma = \{a, b\}$, and let the formal language L be defined as $L = \{w b a b w^T \mid w \in \Sigma^*\}$. Prove or disprove L to be regular.