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## 演習第十

1. Let  $n \in \mathbb{N}^+$  and let  $\mathcal{REG}_n$  be the class of all regular languages that can be accepted by a DFA having at most  $n$  states. Furthermore, let  $\mathcal{CF}_n$  be the class of all context-free languages that can be accepted via empty stack by a PDA having at most  $n$  states.

Prove or disprove the following assertions:

- (a)  $\mathcal{REG}_n \subset \mathcal{REG}_{n+1}$  for all  $n \in \mathbb{N}^+$ ;  
(b)  $\mathcal{CF}_n \subset \mathcal{CF}_{n+1}$  for all  $n \in \mathbb{N}^+$ .

Does your answer to (a) or (b) change if we fix the alphabet  $\Sigma$  to be  $\{a, b\}$ ?

2. Prove or disprove the following.

The language  $L = \{a^n b^m \mid n, m \in \mathbb{N}^+ \text{ and } 1 \leq m \leq n^2\}$  is context-free.

3. Consider the following grammar  $\mathcal{G} = [\{a, b\}, \{\sigma, B, K, S, W\}, \sigma, P]$ , where

$$P = \{\sigma \rightarrow SaK, aK \rightarrow WbbK, aW \rightarrow Wbb, SWb \rightarrow SaB, SWb \rightarrow aB, Bb \rightarrow aB, BK \rightarrow K, BK \rightarrow \lambda\}.$$

(3.1) Determine  $L(\mathcal{G})$ .

(3.2) Prove the correctness of the assertion you made in (3.1).

(3.3) Prove or disprove  $\mathcal{G}$  to be context-sensitive.

4. **Bonus problem:** Let the following pushdown automaton

$\mathcal{K} = [\{q_0, q_1, q_2\}, \{0, 1\}, \{0, 1\}, \delta, q_0, 1, \{q_2\}]$  be given, where

$$\delta(q_i, x, y) = \begin{cases} (q_0, xy), & \text{if } i = 0, x = 0, y \in \{0, 1\}; \\ \{(q_1, \lambda), (q_2, \lambda)\}, & \text{if } i \in \{0, 1\}, x = 1, y \in \{0, 1\}. \end{cases}$$

Provide a context-free grammar  $\mathcal{G}$  for  $L(\mathcal{K})$  by using the construction given in the proof of Theorem 9.3.